# Lecture 7. VC Theory

**COMP90051 Statistical Machine Learning** 

Lecturer: Feng Liu



#### This lecture

- PAC learning bounds:
  - Countably infinite case works as we've done so far
  - \* General infinite case? Needs new ideas!
- Growth functions for the general PAC case
  - \* Considering patterns of labels possible on a data set
  - Gives good PAC bounds provided possible patterns don't grow too fast in the data set size
- Vapnik-Chervonenkis (VC) dimension
  - Max number of points that can be labelled in all ways
  - \* Beyond this point, growth function is polynomial in data set size
  - Leads to famous, general PAC bound from VC theory
- Optional proofs at end (just for fun)

# Countably infinite $\mathcal{F}$ ?

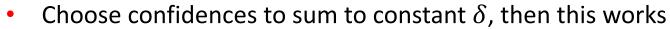
• Hoeffding gave us for a single  $f \in \mathcal{F}$ 

$$\Pr\left(R[f] - \widehat{R}[f] \ge \sqrt{\frac{\log\left(\frac{1}{\delta(f)}\right)}{2m}}\right) \le \delta(f)$$

...where we're free to choose (varying)  $\delta(f)$  in [0,1].

Union bound "works" (sort of) for this case

$$\Pr\left(\exists f \in \mathcal{F}, R[f] - \widehat{R}[f] \ge \sqrt{\frac{\log\left(\frac{1}{\delta(f)}\right)}{2m}}\right) \le \sum_{f \in \mathcal{F}} \delta(f)$$



\* E.g. 
$$\delta(f) = \delta \cdot p(f)$$
 where  $1 = \sum_{f \in \mathcal{F}} p(f)$ 



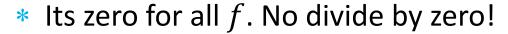
Josh Staiger (CCA2.0)

Try this for finite  ${\mathcal F}$  with uniform p(f)

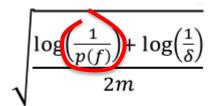
• By inversion: w.h.p 
$$1 - \delta$$
, for all  $f$ ,  $R[f] \le \widehat{R}[f] + \sqrt{\frac{\log(\frac{1}{p(f)}) + \log(\frac{1}{\delta})}{2m}}$ 

# Ok fine, but general case?

- Much of ML has continuous parameters
  - Countably infinite covers only discrete parameters
- - \* p(f) becomes a density



- Need a new argument!
- Idea introduced by VC theory: intuition
  - \* Don't focus on whole class  $\mathcal{F}$  as if each f is different
  - \* Focus on differences over sample  $Z_1, \dots, Z_m$



## Mini Summary

- Can seek out PAC bounds on countably infinite families using Hoeffding bound + union bound
- No good for general (uncountably infinite) cases
- Need another fundamentally new idea

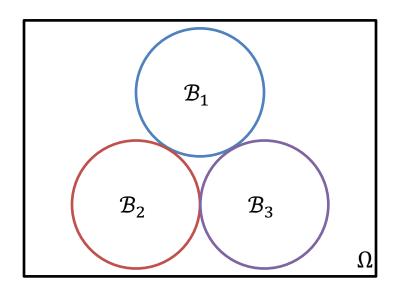
Next: Organising analysis around patterns of labels possible on a data set, to avoid wort-case bad events

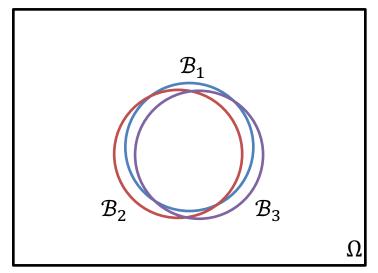
# **Growth Function**

Focusing on the size of model families on data samples

# Bad events: Unreasonably worst case?

- Bad event  $\mathcal{B}_i$  for model  $f_i$   $R[f_i] \hat{R}[f_i] \ge \varepsilon \text{ with probability} \le 2 \exp(-2m\varepsilon^2)$
- Union bound: bad events don't overlap!?  $\Pr(\mathcal{B}_1 \text{ or } ... \text{ or } \mathcal{B}_{|\mathcal{F}|}) \leq \Pr(\mathcal{B}_1) + \cdots + \Pr(\mathcal{B}_{|\mathcal{F}|}) \leq 2|\mathcal{F}| \exp(-2m\varepsilon^2)$

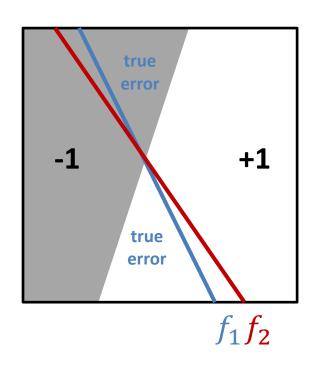




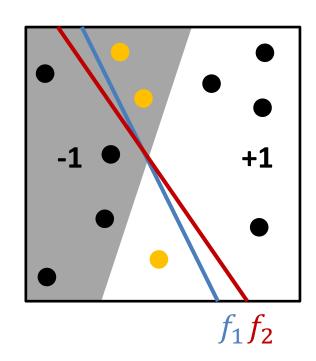
Tight bound: No overlaps

Loose bound: Overlaps

## How do overlaps arise?



Whole of population

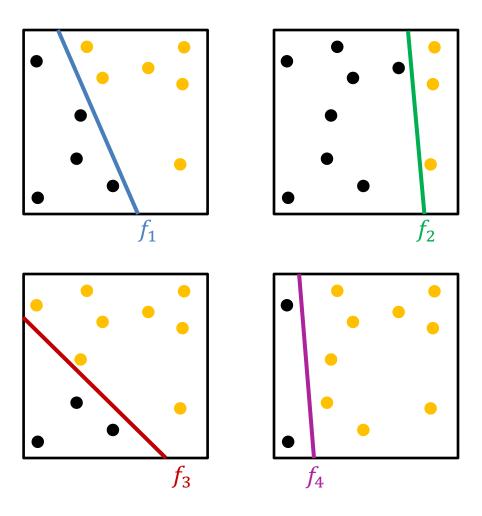


On a sample

Significantly overlapping events  $\mathcal{B}_1$  and  $\mathcal{B}_2$ 

# How do overlaps arise?

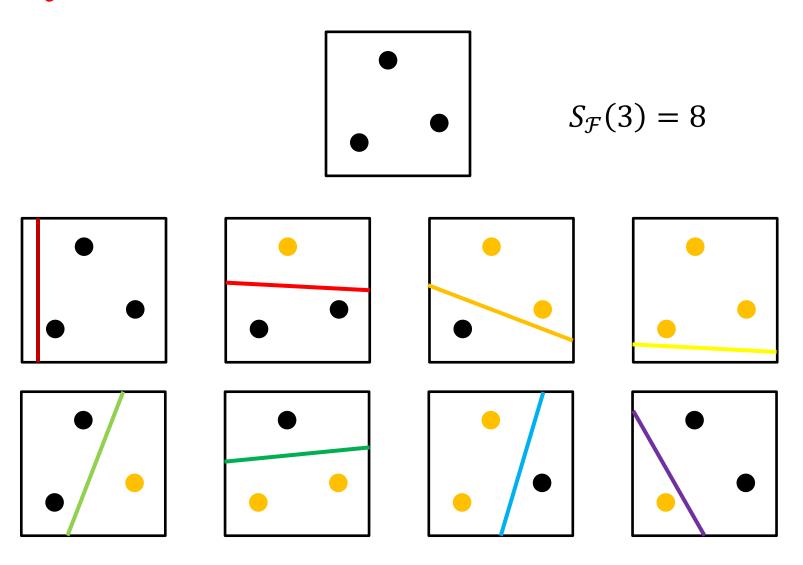
VC theory focuses on the pattern of labels any  $f \in \mathcal{F}$  could make



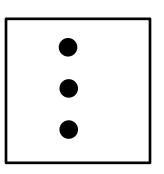
#### Dichotomies and Growth Function

- <u>Definition</u>: Given sample  $x_1, ..., x_m$  and family  $\mathcal{F}$ , a **dichotomy** is a  $(f(x_1), ..., f(x_m)) \in \{-1, +1\}^m$  for some  $f \in \mathcal{F}$ .
- Unique dichotomies  $\mathcal{F}(\mathbf{x}) = \{(f(x_1), ..., f(x_m)) : f \in \mathcal{F}\}$ , patterns of labels possible with the family
- Even when  $\mathcal{F}$  infinite,  $|\mathcal{F}(\mathbf{x})| \leq 2^m$  (why?)
  - And also (relevant for  $\mathcal{F}$  finite, tiny),  $|\mathcal{F}(\mathbf{x})| \leq |\mathcal{F}|$  (why?)
- Intuition:  $|\mathcal{F}(x)|$  might replace  $|\mathcal{F}|$  in union bound? How remove  $\mathbf{x}$ ?
- <u>Definition</u>: The growth function  $S_{\mathcal{F}}(m) = \sup_{\mathbf{x} \in \mathcal{D}^m} |\mathcal{F}(\mathbf{x})|$  is the max number of label patterns achievable by  $\mathcal{F}$  for any m sample.

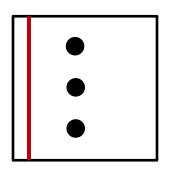
# $S_{\mathcal{F}}(3)$ for $\mathcal{F}$ linear classifiers in 2D?

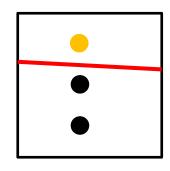


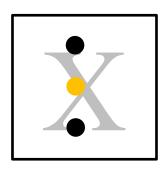
# $S_{\mathcal{F}}(3)$ for $\mathcal{F}$ linear classifiers in 2D?

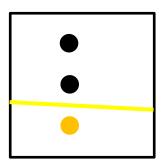


$$|\mathcal{F}(\mathbf{x})| = 6$$
  
but still have  
 $S_{\mathcal{F}}(3) = 8$ 

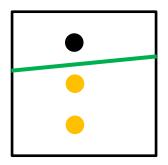


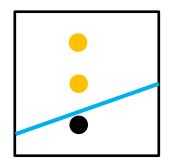


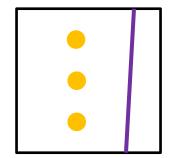






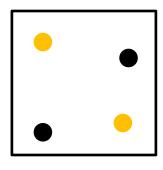


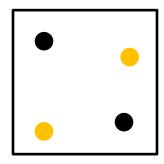




# $S_{\mathcal{F}}(4)$ for $\mathcal{F}$ linear classifiers in 2D?

- What about m = 4 points?
- Can never produce the criss-cross (XOR) dichotomy





- In fact  $S_{\mathcal{F}}(4) = 14 < 2^4$
- Guess/exercise: What about general m and dimension?

#### PAC Bound with Growth Function

• Theorem: Consider any  $\delta>0$  and any class  $\mathcal F$ . Then w.h.p. at least  $1-\delta$ : For all  $f\in\mathcal F$ 

$$R[f] \le \widehat{R}[f] + 2\sqrt{2\frac{\log S_{\mathcal{F}}(2m) + \log(4/\delta)}{m}}$$

- Proof: out of scope ("only" 2-3pgs), optional reading.
- Compare to PAC bounds so far
  - \* A few negligible extra constants (the 2s, the 4)
  - \*  $|\mathcal{F}|$  has become  $S_{\mathcal{F}}(2m)$
  - \*  $S_{\mathcal{F}}(m) \leq |\mathcal{F}|$ , not "worse" than union bound for finite  $\mathcal{F}$
  - \*  $S_{\mathcal{F}}(m) \leq 2^m$ , very bad for big family with exponential growth function gets  $R[f] \leq \hat{R}[f] + Big\ Constant$ . Even  $R[f] \leq \hat{R}[f] + 1$  meaningless!!

# Mini Summary

- The previous PAC bound approach that organises bad events by model and applies uniform bound is only tight if bad events are disjoint
- In reality some models generate overlapping bad events
- Better to organise families by possible patterns of labels on a data set: the dichotomies of the family
- Counting possible dichotomies gives the growth function
- PAC bound with growth function potentially tackles general (uncountably infinite) families provided growth function is sub-exponential in data size

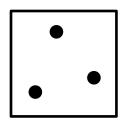
Next: VC dimension for a computable bound on growth functions, with the polynomial behaviour we need! Gives our final, general, PAC bound

# The VC dimension

Computable, bounds growth function

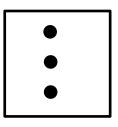
# Vapnik-Chervonenkis dimension

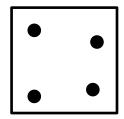
- <u>Definition</u>: The VC dimension VC( $\mathcal{F}$ ) of a family  $\mathcal{F}$  is the largest m such that  $S_{\mathcal{F}}(m)=2^m$ .
  - \* Points  $\mathbf{x} = (x_1, ..., x_m)$  are shattered by  $\mathcal{F}$  if  $|\mathcal{F}(\mathbf{x})| = 2^m$
  - \* So  $VC(\mathcal{F})$  is the size of the largest set shattered by  $\mathcal{F}$
- Example: linear classifiers in  $\mathbb{R}^2$ ,  $VC(\mathcal{F}) = 3$



**Shattered** 

Not shattered





• Guess: VC-dim of linear classifiers in  $\mathbb{R}^d$ ?

## Example: $VC(\mathcal{F})$ from $\mathcal{F}(\mathbf{x})$ on whole domain?

	24	24	24
$\underline{x_1}$	$x_2$	$x_3$	$\frac{x_4}{}$
0	0	0	0
0	1	1	0
1	0	0	1
1	1	0	1
0	1	0	0
1	0	1	0
1	1	1	1
0	0	1	1
0	1	0	1
1	1	1	0

Note we're using labels {0,1} instead of {-1,+1}. Why OK?

- Columns are all points in domain
- Each row is a dichotomy on entire input domain
- Obtain dichotomies on a subset of points  $\mathbf{x}' \subseteq \{x_1, ..., x_4\}$  by: drop columns, drop dupe rows
- $\mathcal{F}$  shatters  $\mathbf{x}'$  if number of rows is  $2^{|\mathbf{x}'|}$

$\overline{x_1}$	$x_2$	$x_4$
$\frac{x_1}{0}$	$\frac{x_2}{0}$	$\frac{x_4}{0}$
0	1	0
1	0	1
1	1	1
0	1	0
1	0	0
1	1	1
0	0	1
0	1	1
1	1	0

#### This example:

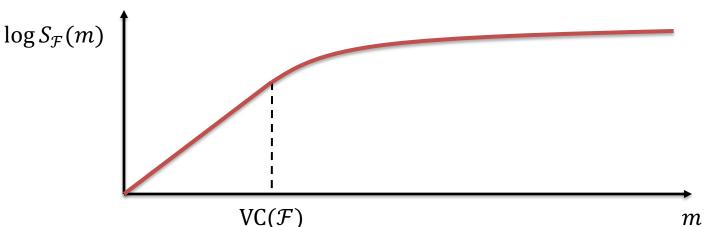
- Dropping column 3 leaves 8 rows behind:  $\mathcal{F}$  shatters  $\{x_1, x_2, x_4\}$
- Original table has  $< 2^4$  rows:  $\mathcal{F}$  doesn't shatter more than 3
- $VC(\mathcal{F}) = 3$

#### Sauer-Shelah Lemma

- Lemma (Sauer-Shelah): Consider any  $\mathcal{F}$  with finite  $VC(\mathcal{F}) = k$ , any sample size m. Then  $S_{\mathcal{F}}(m) \leq \sum_{i=0}^k {m \choose i}$ .
- From basic facts of Binomial coefficients
  - \* Bound is  $O(m^k)$ : finite VC  $\Rightarrow$  eventually polynomial growth!
  - \* For  $m \ge k$ , it is bounded by  $\left(\frac{em}{k}\right)^k$
- Theorem (VC bound): Consider any  $\delta > 0$  and any VC-k class  $\mathcal{F}$ . Then w.h.p. at least  $1 \delta$ : For all  $f \in \mathcal{F}$

$$R[f] \le \widehat{R}[f] + 2\sqrt{2\frac{k\log\frac{2em}{k} + \log\frac{4}{\delta}}{m}}$$





- (Uniform) difference between R[f],  $\widehat{R}[f]$  is  $O\left(\sqrt{\frac{k \log m}{m}}\right)$  down from  $\infty$
- Limiting complexity of  ${\mathcal F}$  leads to better generalisation
- VC dim, growth function measure "effective" size of  ${\cal F}$
- VC dim doesn't count functions, but uses geometry of family: projections of family members onto possible samples
- Example: linear "gap-tolerant" classifiers (like SVMs) with "margin"  $\Delta$  have VC =  $O(1/\Delta^2)$ . Maximising "margin" reduces VC-dimension.

## Mini Summary

- VC-dim is the largest set size shattered by a family
  - \* It is d+1 for linear classifiers in  $\mathbb{R}^d$
  - Can calculate it on entire-domain dichotomies of a family by dropping columns and counting unique rows
- Sauer-Shelah: The growth function grows only polynomially in the set size beyond the VC-dim
- As a result, VC PAC bounds uniform risk and empirical risk deviation by  $O(\sqrt{(VC(\mathcal{F})\log m)/m})$

Next: Two selected proofs. Optional but beautiful.

# **Two Selected Proofs**

Green slides: Not examinable.

Food for thought. Soul food.

### Linear classifiers in d-dim: $VC(\mathcal{F}) \ge d + 1$

• Goal: construct m=d+1 specific points in  $\mathbb{R}^d$  that are shattered by the linear classifier family

• Data in rows of 
$$\mathbf{X} = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 \\ 1 & 1 & 0 & \dots & 0 \\ 1 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & 0 & 0 & \cdots & 1 \end{pmatrix}$$
 is invertible!

- Any dichotomy  $y \in \{-1,1\}^{d+1}$ , need **w** with sign(**Xw**) = **y**
- $\mathbf{w} = \mathbf{X}^{-1}\mathbf{y} \text{ works}!! \operatorname{sign}(\mathbf{X}\mathbf{w}) = \operatorname{sign}(\mathbf{X}\mathbf{X}^{-1}\mathbf{y}) = \operatorname{sign}(\mathbf{y}) = \mathbf{y}$
- We've shown that  $\mathcal{F}$  can shatter d+1 points:  $VC(\mathcal{F}) \geq d+1$

### Linear classifiers in d-dim: $VC(\mathcal{F}) \leq d+1$

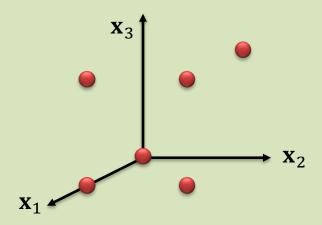
- Goal: cannot shatter any set of d + 2 points
- Any  $\mathbf{x}_1, \dots, \mathbf{x}_{d+2}$ , have more pts than dims: linear dependent  $\mathbf{x}_j = \sum_{i \neq j} a_i \mathbf{x}_i$ , for some j, where not all  $a_i$ 's are zero
- Possible dichotomy **y**?  $y_i = \begin{cases} sign(a_i), & \text{if } i \neq j \\ -1, & \text{if } i = j \end{cases}$ 
  - \* Suppose **w** generated  $i \neq j$ :  $sign(a_i) = sign(\mathbf{w}'\mathbf{x}_i)$  so  $a_i\mathbf{w}'\mathbf{x}_i > 0$
  - \* Can w generate i = j??
  - \*  $\mathbf{w}'\mathbf{x}_j = \mathbf{w}' \sum_{i \neq j} a_i \mathbf{x}_i = \sum_{i \neq j} a_i \mathbf{w}' \mathbf{x}_i > 0$  so  $\operatorname{sign}(\mathbf{w}' \mathbf{x}_j) \neq y_i$
- We've shown  $VC(\mathcal{F}) < d + 2$ , in other words  $VC(\mathcal{F}) = d + 1$

### Proof of Sauer-Shelah Lemma (by Haussler '95)

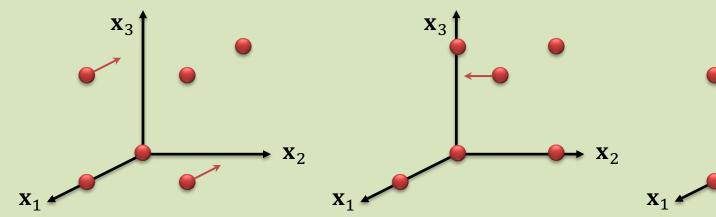
- To show that growth function  $S_{\mathcal{F}}(m) \leq \sum_{i=0}^k {m \choose i}$  we prove the bound for any dichotomies  $|\mathcal{F}(x_1, ..., x_m)|$  since  $|\mathcal{F}(x_1, ..., x_m)| \leq S_{\mathcal{F}}(m)$
- Write  $\mathbf{Y} = \mathcal{F}(x_1, ..., x_m) \subseteq \{0,1\}^m$ , where  $-1 \to 0$ .
- <u>Definition</u>: Consider any column  $1 \le i \le m$  and dichotomy  $\mathbf{y} \in \mathbf{Y}$ . The shift operator  $H_i(\mathbf{y}; \mathbf{Y})$  returns  $\mathbf{y}$  if there exists some  $\mathbf{y}' \in \mathbf{Y}$  differing to  $\mathbf{y}$  only in the  $i^{\text{th}}$  coordinate; otherwise it returns  $\mathbf{y}$  with  $y_i = 0$ . Define  $H_i(\mathbf{Y}) = \{H_i(\mathbf{y}; \mathbf{Y}) : \mathbf{y} \in \mathbf{Y}\}$  the shifting all dichotomies.
  - \* Intuition: Shifting along a column drops a +1 to 0 in that column so long as now other row would become duplicated.
- <u>Definition</u>: A set of dichotomies  $\mathbf{V} \subseteq \{0,1\}^m$  is called closed below if for all  $1 \le i \le m$ , shifting does nothing  $H_i(\mathbf{V}) = \mathbf{V}$ .
  - \* Intuition: Every  $\mathbf{v} \in \mathbf{V}$  has, for every  $1 \le i \le m$  for which  $v_i = 1$ , some  $\mathbf{u} \in \mathbf{V}$  the same as  $\mathbf{v}$  except with  $u_i = 0$ .

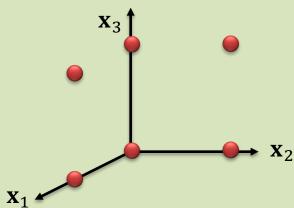
### Proof of Sauer-Shelah Lemma (by Haussler '95)

$x_1$	$x_2$	<i>x</i> <sub>3</sub>
0	0	0
0	1	1
1	0	0
1	1	0
1	0	1
1	1	1



• Example set of 6 unique dichotomies on m=3 pts with VC=2





Shift down along i = 1

Shift down along i = 2

Closed below

### Proof of Sauer-Shelah Lemma (by Haussler '95)

- Goal: show that (1) shifting almost maintains VC dimension and cardinality all the way to a closed-below end, (2) closed-below sets have the desired Sauer-Shelah bound
- Shifting property 1:  $|H_i(\mathbf{Y})| = |\mathbf{Y}|$  for any  $\mathbf{Y}$ .
  - \* Proof: no two dichotomies in Y shift to the same dichotomy
- Shifting property 2:  $VC(H_i(Y)) \leq VC(Y)$  for any i, Y.
  - \* Proof sketch: If  $H_i(\mathbf{Y})$  shatters a subset of points, then so too does  $\mathbf{Y}$
- Shifting property 3: if Y is closed below, then all dichotomies  $y \in Y$  have at most VC(Y)-many  $y_i = 1$  (the rest 0).
  - \* Therefore:  $|\mathbf{Y}| \leq {m \choose 0} + {m \choose 1} + \dots + {m \choose VC(\mathbf{Y})}$  by counting
  - \* Proof sketch: if a  $y \in Y$  had more 1s, all combinations would exist "below"
- Together: exists a shift sequence  $i_1, \dots, i_N$  to a closed below  $H_{i_N}(\mathbf{Y})$ :  $|\mathbf{Y}| = \left| H_{i_1}(\mathbf{Y}) \right| = \dots = \left| H_{i_N}(\mathbf{Y}) \right| \leq \sum_{i=0}^{\text{VC}\left(H_{i_N}(\mathbf{Y})\right)} {m \choose i} \leq \dots \leq \sum_{i=0}^{\text{VC}(\mathbf{Y})} {m \choose i}$

### Mini Summary

- Linear classifiers in  $\mathbb{R}^d$  have VC dimension d+1
  - Lower bound VC-dim with specific points that are shattered
  - \* Upper bound VC-dim by lin. dependence of any d+2 points
- Sauer-Shelah lemma bounds a family's growth function by a polynomial in VC dimension.
  - Ingenious shifting operator transforms sets of dichotomies into boundable closed-below sets
  - \* Along the way keeps cardinality and VC-dim controlled

Next time: Support vector machines